

# Causal Graphical Models (Summer Short Course)

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### 0.1 Literature

#### 0.1.1 Day 1 (Lecture 1-4)

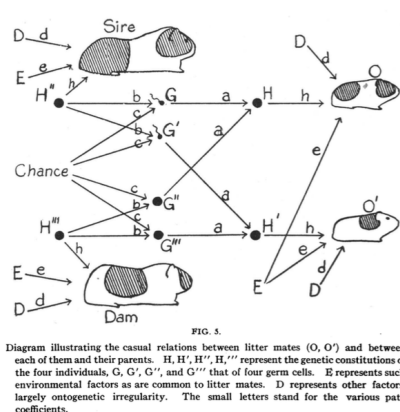
- Zhao, Q. (2024). A matrix algebra for graphical statistical models. (arXiv:2407.15744)

#### 0.1.2 Day 2 (Lecture 5-8)

- Lauritzen, S. L. (1996). *Graphical Models*. Clarendon Press.
- Zhao, Q. (2025). On statistical and causal models associated with acyclic directed mixed graphs. (arXiv:2501.03048)
- Guo, F. R., & Zhao, Q. (2023). Confounder selection via iterative graph expansion. (arXiv:2309.06053)

# 1 Lecture 1: Directed mixed graphs and linear systems

## 1.1 History



- **Sewall Green Wright** (December 21, 1889 – March 3, 1988) was an American geneticist known for his influential work on evolutionary theory and also for his work on **path analysis**. He was a founder of population genetics alongside Ronald Fisher and J. B. S. Haldane, which was a major step in the development of the modern synthesis combining genetics with evolution.

## 1.2 Directed mixed graphs

We will consider graphs with two types of edges: directed ( $\rightarrow$ ) and bidirected ( $\leftrightarrow$ ).

### 1.2.1 Definition

A **directed mixed graph (DMG)**  $G = (V, \mathcal{D}, \mathcal{B})$  consists of a finite vertex set  $V$ , a directed edge set  $\mathcal{D} \subseteq V \times V$  that contains ordered pairs of vertices, and a bidirected edge set  $\mathcal{B} \subseteq V \times V$  that contains unordered pairs of vertices (so  $(j, k) \in \mathcal{B}$  implies  $(k, j) \in \mathcal{B}$ ) such that

$$(j, k) \in \mathcal{B} \implies (j, j) \in \mathcal{B}, (k, k) \in \mathcal{B}, \quad \text{for all } j, k \in V.$$

Let  $\mathbb{G}(V)$  denote the collection of all directed mixed graphs with vertex set  $V$ .

We say the **directed edge** “ $j \rightarrow k$ ” is contained in  $G$  if  $(j, k) \in \mathcal{D}$ , and in this case we say this is an *incoming edge* for  $k$ , an *outgoing edge* for  $j$ , the vertex  $j$  is a **parent** of  $k$ , and  $k$  is a **child** of  $j$  in  $G$ . Likewise, we say the **bidirected edge** “ $j \leftrightarrow k$ ” is contained in  $G$  if  $(j, k) \in \mathcal{B}$ .

## 1.3 Causal interpretation

- Directed edges mean direct causal effects.
- Bidirected edges mean unspecified, residual/exogenous correlations.

### 1.3.1 Why directed edges?

- Causality is **transitive** ( $A$  causes  $B$  and  $B$  causes  $C \implies A$  causes  $C$ ). This defines a **pre-order**.
  - This can be described by the reachability relationship of a **directed graph**.
- Often we think causality is **irreflexive** ( $A$  does not cause itself) and **asymmetric** ( $A$  and  $B$  cannot be causes of each other). This defines a **partial order**.
  - This can be described by the reachability relationship of a **directed acyclic graph**.

### 1.3.2 Why bidirected edges?

- In statistics and causal inference we are often concerned with **latent variables**.

## 1.4 Canonical graphs

### 1.4.1 Definition

- We say the directed mixed graph is **canonical** if it contains all bidirected loops.
  - The full collection with vertex set  $V$  is denoted by  $\mathbb{G}^*(V)$ .
- We say the graph is **canonically directed** if it is canonical and contains no other bidirected edges.
  - The full collection with vertex set  $V$  is denoted by  $\mathbb{G}_D^*(V)$ .

For such graphs, it is usually more convenient to use the trimmed graph obtained by

$$\text{trim} : (V, \mathcal{D}, \mathcal{B}) \mapsto (V, \mathcal{D}, \mathcal{B} \setminus \{(j, j) : j \in V\}).$$

## 1.5 Walk, path, cycle

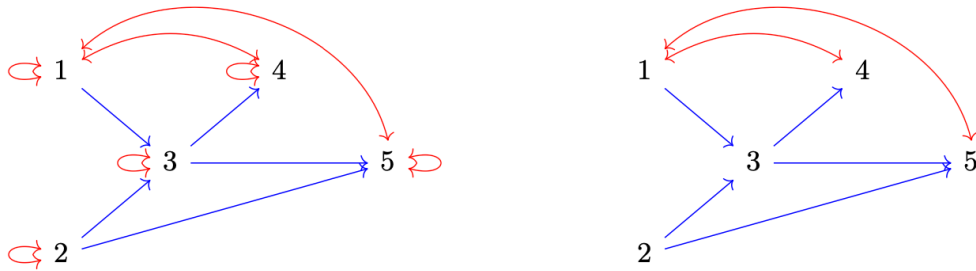
### 1.5.1 Definition

- A **walk** is an ordered sequence of connected edges ignoring edge direction.
- A **path** is a walk with no repeated vertices.
- A **cycle** is a walk with the same starting and ending vertices.
- Vertices at the two ends of a walk are called its **endpoints**, and the other vertices are called **non-endpoints**.
- A walk is **directed** if all its edges have the same direction, like  $j \longrightarrow \dots \longrightarrow k$ .
- The graph  $G$  (or its direct subgraph) is **acyclic** if it contains no directed cycles.

### 1.5.2 Notation

- $\mathbb{G}_A^*(V)$ : the collection of acyclic, canonical DMGs with vertex set  $V$ .
- $\mathbb{G}_{DA}^*(V)$ : the collection acyclic, canonically directed DMGs (basically DAGs).

## 1.6 Demonstration of trimming



Is the directed mixed graph on the left

- canonical?
- canonically directed?
- acyclic?

## 1.7 Gaussian linear system on a graph

We will not distinguish a random vector  $V = (V_1, \dots, V_d)$  (probability theory) with a vertex set of random variables  $V = \{V_1, \dots, V_d\}$  (graph theory).

### 1.7.1 Definition

We say a random vector  $V$  follows a **Gaussian linear system** wrt  $G \in \mathbb{G}(V)$  if

$$V = \beta^T V + E \text{ for some } E \sim N(0, \Lambda)$$

where

- $\beta \in \mathbb{R}^{d \times d}$  respects the directed subgraph:  $V_j \not\rightarrow V_k$  in  $G \Rightarrow \beta_{jk} = 0$ .
- $\Lambda \in \mathbb{R}^{d \times d}$  respects the bidirected subgraph:  $V_j \not\leftrightarrow V_k$  in  $G \Rightarrow \Lambda_{jk} = 0$ .

### 1.7.2 Gaussian model

- Let  $\mathbb{N}(G)$  denote the collection of probability distributions of such  $V$ .
- Let  $\mathbb{N}^+(G)$  denote the subclass where  $\Lambda$  is positive definite (so  $G \in \mathbb{G}^*(V)$ ).

## 1.8 Roadmap: Basic questions

Let  $V$  be a random vector and  $J = V_{\mathcal{J}}$ ,  $K = V_{\mathcal{K}}$ ,  $L = V_{\mathcal{L}}$  be sub-vectors of  $V$ .

1. What is the probability distribution of  $J$ ?
2. Is  $J \perp\!\!\!\perp K$  true?
3. Is  $J \perp\!\!\!\perp K \mid L$  true?

### 1.8.1 Answer

If  $V \sim N(0, \Sigma)$  and  $\Sigma$  is positive definite, then

1.  $J \sim N(0, \Sigma_{\mathcal{J}, \mathcal{J}})$ .
2.  $J \perp\!\!\!\perp K$  if and only if  $\Sigma_{\mathcal{J}, \mathcal{K}} = 0$ .
3.  $V_j \perp\!\!\!\perp V_k \mid V_{[d] \setminus \{j, k\}}$  if and only if  $(\Sigma^{-1})_{jk} = 0$ .

### 1.8.2 Lecture 1-3

Let  $P \in \mathbb{N}(G)$  be the distribution of  $V$ . Can we answer these questions using just  $G$ ?

## 1.9 Roadmap: From linear algebra to graphs

### 1.9.1 Basic combinatoric result

Let  $A$  be the adjacency matrix of a directed graph  $G$ . Then

$$(A^r)_{jk} = |\{\text{directed walks from } j \text{ to } k \text{ of length } r\}|, \quad r \geq 1,$$

$$[(Id - A)^{-1}]_{jk} = [Id + A + A^2 + \dots]_{jk} = \delta_{jk} + |\{\text{directed walks from } j \text{ to } k\}|.$$

So **matrix multiplication is similar to walking on a graph**.

- Thinking abstractly, edges in a graph encode certain local relations. By composing ("multiplying") those edges, we can obtain new global relations.

## 1.10 Matrices of walks

### 1.10.1 Basic matrices

Edges in  $G = (V, \mathcal{D}, \mathcal{B}) \in \mathbb{G}(V)$  can be rearranged into:

$$W[j \longrightarrow k \text{ in } G] = \begin{cases} \{j \longrightarrow k\}, & \text{if } (j, k) \in \mathcal{D}, \\ \emptyset, & \text{otherwise,} \end{cases}$$

$$W[j \longleftrightarrow k \text{ in } G] = \begin{cases} \{j \longleftrightarrow k\}, & \text{if } (j, k) \in \mathcal{B}, \\ \emptyset, & \text{otherwise.} \end{cases}$$

- Examples (on blackboard).

## 1.11 Basic operations

### 1.11.1 On sets of walks

1. Addition  $+$  means set union.

- Example:  $\{V_2 \longrightarrow V_5\} + \{V_2 \longrightarrow V_3 \longrightarrow V_5\} = ?$ .

2. Multiplication  $\cdot$  means concatenation.

- Example:  $\{V_2 \longleftrightarrow V_2\} \cdot \{V_2 \longrightarrow V_5, V_2 \longrightarrow V_3 \longrightarrow V_5\} = ?$ .

3. Transpose  $^T$  means reversing direction.

- Example:  $\{V_2 \longrightarrow V_5, V_2 \longrightarrow V_3 \longrightarrow V_5\}^T = ?$ .

### 1.11.2 On matrices

Examples (on blackboard).

$$(W + W')[V_j, V_k] = W[V_j, V_k] + W'[V_j, V_k],$$

$$(W \cdot W')[V_j, V_k] = \sum_{V_l \in V} W[V_j, V_l] \cdot W'[V_l, V_k],$$

$$(W^T)[V_j, V_k] = (W[V_k, V_j])^T = \{w^T : w \in W[V_k, V_j]\}.$$

## 1.12 Further definitions

- **Right-directed walks:**  $W[V \rightsquigarrow V] = \sum_{q=1}^{\infty} (W[V \longrightarrow V])^q$ .
- **Left-directed walks:**  $W[V \leftrightsquigarrow V] = (W[V \rightsquigarrow V])^T$ .
- **Identity matrix** (for multiplication):  $\text{Id} = \text{diag}(\text{id}, \dots, \text{id})$ , where  $\text{id}$  is the trivial walk with length 0 such that

$$\text{id} \cdot w = w \cdot \text{id} = w \quad \text{and} \quad \text{Id} \cdot W = W \cdot \text{Id} = W$$

- Let  $\emptyset$  also denote the matrix with empty sets of walks. This is the identity element for matrix addition (set union).

### 1.13 Weight function

- Recall that  $P \in \mathbb{N}(G)$  means  $V = \beta^T V + E$  where  $E \sim N(0, \Lambda)$ .
- So  $V = (Id - \beta)^{-T} E \sim N(0, \Sigma)$  where

$$\Sigma = (Id - \beta)^{-T} \Lambda (Id - \beta)^{-1}.$$

How can we represent this graphically?

- Let  $\sigma$  be the **weight function** on all walks in  $G$  generated by

$$\beta = \sigma(W[V \rightarrow V]) \quad \text{and} \quad \Lambda = \sigma(W[V \leftrightarrow V]).$$

- Example:  $\sigma(\{V_1 \leftrightarrow V_4, V_1 \leftrightarrow V_1 \rightarrow V_3 \rightarrow V_4\}) = \Lambda_{14} + \Lambda_{11}\beta_{13}\beta_{34}$ .

#### 1.13.1 Lemma

If  $\beta$  is stable (spectral radius  $< 1$ ), then

$$(Id - \beta)^{-1} = Id + \sigma(W[V \rightsquigarrow V]).$$

## 2 Lecture 2: Path analysis and graph marginalization

### 2.1 Trek rule

This motivates us to define **treks** or **t-connected walks** in  $G$  as (expand on blackboard)

$$W[V \overset{t}{\longleftrightarrow} V] = (Id + W[V \rightsquigarrow V]) \cdot W[V \leftrightarrow V] \cdot (Id + W[V \rightsquigarrow V]).$$

#### 2.1.1 Theorem

Suppose  $G \in \mathbb{G}(V)$  and  $P \in \mathbb{N}(G)$  with weight function  $\sigma$ , then

$$\text{Cov}_P(V) = \sigma(W[V \overset{t}{\longleftrightarrow} V \text{ in } G]).$$

#### 2.1.2 Examples

- $\text{Var}(V_3) = ?$
- $\text{Cov}(V_3, V_4) = ?$

### 2.2 Arcs

#### 2.2.1 Definitions

- We say a walk is an **arc** or **m-connected** if it has no collider (like  $\rightarrow j \leftarrow$ ).
- We say a walk is **d-connected** if it is an arc and contains no bidirected edge, so

#### 2.2.2 Proposition

An arc has exactly zero or one bidirected edge.

#### 2.2.3 Notation

A squiggly line ( $\rightsquigarrow$ ) means no collider, and we use no/half/full arrowheads at both ends.

$$\begin{aligned} W[V \overset{d}{\rightsquigarrow} V] &= W[V \rightsquigarrow V] + W[V \rightsquigarrow V] + W[V \rightsquigarrow V \rightsquigarrow V], \\ W[V \rightsquigarrow V] &= W[V \overset{t}{\longleftrightarrow} V] + W[V \overset{d}{\rightsquigarrow} V]. \end{aligned}$$

## 2.3 From treks to paths

- Notation:  $P[\dots] = W[\dots] \cap \mathcal{P}_G$  where  $\mathcal{P}_G$  contains all paths on  $G$  and  $\cap$  is applied entry-wise.

### 2.3.1 Lemma

For any  $G \in \mathbb{G}^*(V)$  and any  $j, k \in V$ ,  $j \neq k$ , we have

$$j \overset{t}{\rightsquigarrow} k \text{ in } G \iff P[j \rightsquigarrow k \text{ in } G] \neq \emptyset \iff P[j \rightsquigarrow k \text{ in } \text{trim}(G)] \neq \emptyset.$$

- Proof on blackboard (assuming  $G$  is acyclic).
- Key definition:

$$P[j \overset{d}{\rightsquigarrow} k \text{ via root } r] = \begin{cases} P[j \rightsquigarrow k], & \text{if } r = j, \\ P[j \rightsquigarrow k], & \text{if } r = k, \\ (P[j \rightsquigarrow r] \cdot P[r \rightsquigarrow k]) \cap \mathcal{P}_G, & \text{otherwise.} \end{cases}$$

## 2.4 Wright's path analysis

### 2.4.1 Theorem

Suppose  $G \in \mathbb{G}_A(V)$  and  $P \in \mathbb{N}(G)$  with weight function  $\sigma$ , then for any  $V_j, V_k \in V$ ,  $j \neq k$ , we have

$$\begin{aligned} \text{Cov}_P(V_j, V_k) &= \sigma(P[V_j \overset{t}{\rightsquigarrow} V_k \text{ in } G]) \\ &+ \sum_{V_r \in V} \sigma(P[V_j \overset{d}{\rightsquigarrow} V_k \text{ via root } V_r \text{ in } G]) \cdot \text{Var}_P(V_r). \end{aligned}$$

- Proof on blackboard.
- Example:  $\text{Cov}(V_3, V_4) = ?$

## 2.5 Blocking arcs

### 2.5.1 Definition

We say an arc is **blocked** by  $L \subseteq V$  if the arc has an non-endpoint in  $L$ .

### 2.5.2 Examples

1. Which of the following are blocked by 3?
  - $1 \longrightarrow 3 \longrightarrow 4$
  - $4 \longleftarrow 3 \longleftrightarrow 3 \longrightarrow 5$
  - $1 \longrightarrow 3 \longleftarrow 2$
2. Define  $W[V \rightsquigarrow V \mid L]$  and  $W[V \overset{t}{\rightsquigarrow} V \mid L]$  using the matrix algebra.

### 2.5.3 Lemma

If  $\beta = \sigma(W[V \longrightarrow V])$  is principally stable, then for any  $L \subseteq V$ , we have

$$(\text{Id} - \beta_{L^c L^c})^{-1} = \text{Id} + \sigma(W[L^c \rightsquigarrow L^c \mid L]).$$

## 2.6 Marginalization of graphs

- Notation: Write  $W[\dots \text{ in } G] \neq \emptyset$  as  $\dots \text{ in } G$ .
- For  $G \in \mathbb{G}(V)$  and  $\tilde{V} \subseteq V$ , the **marginal graph**  $\text{margin}_{\tilde{V}}(G)$  is obtained by

$$\begin{aligned} j \longrightarrow k \text{ in } \tilde{G} &\iff j \rightsquigarrow k \mid \tilde{V} \text{ in } G, \quad j, k \in \tilde{V}, \\ j \longleftrightarrow k \text{ in } \tilde{G} &\iff j \overset{t}{\rightsquigarrow} k \mid \tilde{V} \text{ in } G, \quad j, k \in \tilde{V}. \end{aligned}$$

- If  $\tilde{V}$  is **ancestral** (meaning  $\tilde{V}$  contains  $\text{an}(\tilde{V}) = \{V_j \in V : V_j \rightsquigarrow \tilde{V}\}$ ), then  $\text{margin}_{\tilde{V}}(G)$  is the subgraph of  $G$  restricted to  $\tilde{V}$ .
- It is often useful to view  $\text{margin}_{\tilde{V}}$  as a map of walks in  $G$  to walks in  $\tilde{G}$ .

### 2.6.1 Proposition: Marginalization preserves unblocked directed walks and treks

For any  $G \in \mathbb{G}(V)$  and  $L \subseteq \tilde{V} \subseteq V$ ,

$$\text{margin}_{\tilde{V}} \left( W \left[ \tilde{V} \left\{ \begin{array}{c} \rightsquigarrow \\ \rightsquigarrow \\ \overset{t}{\rightsquigarrow} \end{array} \right\} \tilde{V} \mid L \text{ in } G \right] \right) = W \left[ \tilde{V} \left\{ \begin{array}{c} \rightsquigarrow \\ \rightsquigarrow \\ \overset{t}{\rightsquigarrow} \end{array} \right\} \tilde{V} \mid L \text{ in } \text{margin}_{\tilde{V}}(G) \right].$$

## 2.7 Marginalization of linear systems

### 2.7.1 Theorem

For  $G \in \mathbb{G}(V)$  and  $P \in \mathbb{N}(G)$ , if  $\beta$  is principally stable, then

$$\text{margin}_{\tilde{V}}(P) \in \mathbb{N}(\text{margin}_{\tilde{V}}(G))$$

with weight function generated by

$$\begin{aligned} \tilde{\sigma}(W[\tilde{V} \longrightarrow \tilde{V} \text{ in } \tilde{G}]) &= \sigma(W[\tilde{V} \rightsquigarrow \tilde{V} \mid \tilde{V} \text{ in } G]), \\ \tilde{\sigma}(W[\tilde{V} \longleftrightarrow \tilde{V} \text{ in } \tilde{G}]) &= \sigma(W[\tilde{V} \overset{t}{\rightsquigarrow} \tilde{V} \mid \tilde{V} \text{ in } G]). \end{aligned}$$

- Proof of blackboard.

### 2.7.2 Example

- Find the marginal graph and linear system with  $\tilde{V} = \{V_1, V_2, V_4, V_5\}$ .

## 2.8 Marginalization of canonical graphs

### 2.8.1 Definition (marginalization of trimmed graphs)

Consider  $G \in \mathbb{G}^*(V)$  and  $G^* = \text{trim}(G)$ . For  $\tilde{V} \subseteq V$ , the marginal trimmed graph  $\tilde{G}^* = \text{margin}_{\tilde{V}}^*(G^*)$  is obtained by

$$\begin{aligned} j \longrightarrow k \text{ in } \tilde{G}^* &\iff P[j \rightsquigarrow k \mid \tilde{V} \text{ in } G^*] \neq \emptyset, \\ j \longleftrightarrow k \text{ in } \tilde{G}^* &\iff P[j \overset{t}{\rightsquigarrow} k \mid \tilde{V} \text{ in } G^*] \neq \emptyset. \end{aligned}$$

- This is the usual definition in the literature and is justified by the following commutative diagram (proof omitted; example on blackboard).

$$\begin{array}{ccc} \mathbb{G}^*(V) & \xrightarrow{\text{trim}} & \text{trim}(\mathbb{G}^*(V)) \\ \text{margin}_{\tilde{V}} \downarrow & & \downarrow \text{margin}_{\tilde{V}}^* \\ \mathbb{G}^*(\tilde{V}) & \xrightarrow{\text{trim}} & \text{trim}(\mathbb{G}^*(\tilde{V})) \end{array} \quad \begin{array}{ccc} G & \xrightarrow{\text{trim}} & G^* \\ \text{margin}_{\tilde{V}} \downarrow & & \downarrow \text{margin}_{\tilde{V}}^* \\ \tilde{G} & \xrightarrow{\text{trim}} & \tilde{G}^* \end{array}$$

### 3 Lecture 3: m-separation and conditional independence

#### 3.1 Conditional independence in multivariate normal

Suppose  $V \sim N(\mu, \Sigma)$  and  $\Sigma$  is positive definite.

##### 3.1.1 Proposition

1. For  $\mathcal{J} \subset [d]$  and  $\mathcal{L} = [d] \setminus \mathcal{J}$ , we have

$$V_{\mathcal{J}} \mid V_{\mathcal{L}} = v_{\mathcal{L}} \sim N(\mu_{\mathcal{J}} + \Sigma_{\mathcal{J}\mathcal{L}}\Sigma_{\mathcal{L}\mathcal{L}}^{-1}(v_{\mathcal{L}} - \mu_{\mathcal{L}}), \Sigma_{\mathcal{J}\mathcal{J}} - \Sigma_{\mathcal{J}\mathcal{L}}\Sigma_{\mathcal{L}\mathcal{L}}^{-1}\Sigma_{\mathcal{L}\mathcal{J}})$$

2. For any  $j, k \in [d]$ ,  $j \neq k$ , we have

$$V_j \perp\!\!\!\perp V_k \mid V_{[d] \setminus \{j, k\}} \iff \text{Cov}(V_j, V_k \mid V_{[d] \setminus \{j, k\}}) = 0 \iff (\Sigma^{-1})_{jk} = 0.$$

##### 3.1.2 Proof sketch

1.  $V_{\mathcal{J}} - \Sigma_{\mathcal{J}\mathcal{L}}\Sigma_{\mathcal{L}\mathcal{L}}^{-1}V_{\mathcal{L}}$  is independent of  $V_{\mathcal{L}}$ .
2. Use  $\mathcal{J} = \{j, k\}$ , then use the block matrix inverse formula.

#### 3.2 Collider-connected walks

Suppose  $P \in \mathbb{N}^+(G)$  and  $(\text{Id} - \beta)$  is non-singular.

- From Lecture 2:  $\Sigma = (\text{Id} - \beta)^{-T} \Lambda (\text{Id} - \beta)^{-1}$ .
- So  $\Sigma^{-1} = (\text{Id} - \beta) \Lambda^{-1} (\text{Id} - \beta)^T$ . More specifically,

$$(\Sigma^{-1})_{jk} = (\Lambda^{-1})_{jk} - \sum_{m \in [d]} \beta_{jm} (\Lambda^{-1})_{mk} - \sum_{l \in [d]} (\Lambda^{-1})_{jl} \beta_{kl} + \sum_{m, l \in [d]} \beta_{jm} (\Lambda^{-1})_{ml} \beta_{kl}$$

- A sufficient condition for  $(\Sigma^{-1})_{jk} = 0$  is that every RHS term vanishes.

##### 3.2.1 Lemma

For any  $V_j \neq V_k$ , we have

$$W[V_j \longleftrightarrow * \longleftrightarrow V_k \text{ in } G] = \emptyset \implies V_j \perp\!\!\!\perp V_k \mid V \setminus \{V_j, V_k\} \text{ under } P.$$

- $V_j \longleftrightarrow * \longleftrightarrow V_k$  means a walk from  $V_j$  to  $V_k$  where **every non-endpoint is a collider**.
- Proof on blackboard. (*Hint*: First assume  $\beta = 0$ .)

#### 3.3 Towards the general case

- Notation: Write  $W[\dots] = \emptyset$  as **not**  $\dots$ .

##### 3.3.1 Goal

We would like to establish a graphical condition for  $J \perp\!\!\!\perp K \mid L$ .

- Because  $V$  is Gaussian, it suffices to consider  $V_j \perp\!\!\!\perp V_k \mid L$  for all  $V_j \in J, V_k \in K$ .
- Using the last Lemma, we have (let  $\tilde{V} = \{V_j, V_k\} \cup L$ )

$$\text{not } V_j \longleftrightarrow * \longleftrightarrow V_k \text{ in } \text{margin}_{\tilde{V}}(G) \implies V_j \perp\!\!\!\perp V_k \mid L \text{ under } P.$$

### 3.3.2 Main problem

Can we conclude **not**  $V_j \longleftrightarrow * \longleftrightarrow V_k$  in  $\text{margin}_{\tilde{V}}(G)$  using  $G$  directly?

## 3.4 General blocking

### 3.4.1 Definition

- We say a walk in  $G \in \mathbb{G}(V)$  is (ancestrally) blocked by  $L \subseteq V$ , if
  1. the walk contains a collider  $m$  such that  $m \notin L$  (and  $m \not\rightsquigarrow L$ ), **or**
  2. the walk contains a non-colliding non-endpoint  $m$  such that  $m \in L$ .
- Let  $W[V \rightsquigarrow * \rightsquigarrow V \mid L \text{ in } G]$  collect all walks that are not blocked by  $L$ .
- We say  $J, K \subseteq V$  are **m-separated** given  $L$  if **not**  $J \rightsquigarrow * \rightsquigarrow K \mid L$  in  $G$ .

### 3.4.2 Remarks

- All walks are separated by 0, 1, or more colliders.
- $W[V \rightsquigarrow * \rightsquigarrow V] \neq W[V \rightsquigarrow * \rightsquigarrow V \mid \emptyset] = W[V \rightsquigarrow V]$ .
- The literature usually uses ancestral blocking with paths (see next Proposition).

### 3.4.3 Example

- Find all m-separations in the running example. (*Hint*: There are 2 in total.)

## 3.5 Graph separation

- **t-separation**: only consider  $\overset{t}{\rightsquigarrow} * \overset{t}{\rightsquigarrow}$  (one or more treks).
- **d-separation**: only consider walks consisting of directed edges only.

### 3.5.1 Lemma

For any  $G \in \mathbb{G}(V)$  and  $J, K, L \subseteq \tilde{V} \subseteq V$ , we have

$$J \overset{t}{\rightsquigarrow} * \overset{t}{\rightsquigarrow} K \mid L \text{ in } \text{margin}_{\tilde{V}}(G) \iff J \overset{t}{\rightsquigarrow} * \overset{t}{\rightsquigarrow} K \mid L \text{ in } G.$$

### 3.5.2 Lemma

Consider  $G \in \mathbb{G}(V)$  and  $\{j\}, \{k\}, L \subseteq V$ . If  $G$  is canonical, then

$$(i) \ j \overset{t}{\rightsquigarrow} * \overset{t}{\rightsquigarrow} k \mid L \iff (ii) \ j \rightsquigarrow * \rightsquigarrow k \mid L \iff (iii) \ P[j \rightsquigarrow * \rightsquigarrow k \mid_a L] \neq \emptyset.$$

Furthermore, if  $G$  is canonically directed, then

$$(i), (ii), (iii) \iff (iv) \ j \overset{d}{\rightsquigarrow} * \overset{d}{\rightsquigarrow} k \mid L \iff (v) \ P[j \overset{d}{\rightsquigarrow} * \overset{d}{\rightsquigarrow} k \mid_a L] \neq \emptyset.$$

## 3.6 General result

### 3.6.1 Theorem

Suppose  $P \in \mathbb{N}^+(G)$  for some  $G \in \mathbb{G}^*(V)$  and  $(\text{Id} - \beta)$  is principally non-singular. Then for all disjoint  $J, K, L \subseteq V$ , we have

$$\text{not } J \rightsquigarrow * \rightsquigarrow K \mid L \text{ in } G \implies J \perp\!\!\!\perp K \mid L \text{ under } P.$$

- Proof on blackboard using the results above.

## 4 Lecture 4: Linear structural equation model and identifiability

### 4.1 Potential outcome

#### 4.1.1 Motivation

We can represent Hooke's law by ( $X$  force,  $Y$  compression distance,  $\beta$  elasticity):

$$X \longrightarrow Y \quad \text{and} \quad Y = \beta X.$$

- But equations have no direction:  $Y = \beta X$  is equivalent to  $X = Y/\beta$ .
- To emphasize "force causes compression", we can write  $Y(x) = \beta x$  for all  $x$ .

#### 4.1.2 Definition

- Let  $V(V_{\mathcal{I}} = v_{\mathcal{I}})$  (often abbreviated as  $V(v_{\mathcal{I}})$ ) denote the **potential outcome** of the entire system under an intervention that sets  $V_{\mathcal{I}}$  to  $v_{\mathcal{I}}$ .
- A **causal model** is a collection of probability distributions  $P$  on the **potential outcomes schedule**  $V(\cdot) = (V(v_{\mathcal{I}}) : \mathcal{I} \in [d])$  such that

$$P(V(v_{\mathcal{I}}, v_{\mathcal{J}}) = V(v_{\mathcal{I}}) \mid V_{\mathcal{J}}(v_{\mathcal{I}}) = v_{\mathcal{J}}) = 1, \text{ for all disjoint } \mathcal{I}, \mathcal{J} \subseteq [d], v \in \mathbb{V}.$$

### 4.2 Linear structural equation model (Linear SEM)

#### 4.2.1 Definition

We say  $V(\cdot)$  follows a **linear SEM** with respect to  $G \in \mathbb{G}(V)$ , if there exist  $\beta$  and  $\Lambda$  compatible with  $G$  such that

$$V_j(v_{\mathcal{I}}) = \sum_{k \in \text{pa}(j) \cap \mathcal{I}} \beta_{kj} v_k + \sum_{k \in \text{pa}(j) \setminus \mathcal{I}} \beta_{kj} V_k(v_{\mathcal{I}}) + E_j, \text{ for all } j \in [d] \text{ and } \mathcal{I} \subseteq [d],$$

for some  $E = (E_1, \dots, E_d)$  with  $\text{Cov}(E) = \Lambda$  and  $\text{pa}_G(j) = \{l \in [d] : V_l \longrightarrow V_j \text{ in } G\}$ .

- In words, every equation still "holds" (thus is "structural") under any intervention.

#### 4.2.2 Example

In the running example, how do the structural equations look like under the intervention  $(V_2, V_3) = (v_2, v_3)$ ?

### 4.3 Single-world intervention graphs

We can rewrite the structural equations for  $V(v_{\mathcal{I}})$  in matrix form:

$$\begin{pmatrix} v_{\mathcal{I}} \\ V_{\mathcal{I}}(v_{\mathcal{I}}) \\ V_{\mathcal{I}^c}(v_{\mathcal{I}}) \end{pmatrix} = \begin{pmatrix} v_{\mathcal{I}} \\ 0 \\ 0 \end{pmatrix} + \begin{pmatrix} 0 & 0 & 0 \\ \beta_{\mathcal{I}, \mathcal{I}}^T & 0 & \beta_{\mathcal{I}^c, \mathcal{I}}^T \\ \beta_{\mathcal{I}, \mathcal{I}^c}^T & 0 & \beta_{\mathcal{I}^c, \mathcal{I}^c}^T \end{pmatrix} \begin{pmatrix} v_{\mathcal{I}} \\ V_{\mathcal{I}}(v_{\mathcal{I}}) \\ V_{\mathcal{I}^c}(v_{\mathcal{I}}) \end{pmatrix} + \begin{pmatrix} 0 \\ E_{\mathcal{I}} \\ E_{\mathcal{I}^c} \end{pmatrix}.$$

So  $(v_{\mathcal{I}}, V(v_{\mathcal{I}}))$  follows a linear system with respect to the **single-world intervention graph (SWIG)**  $G(v_{\mathcal{I}})$  obtained by modifying  $G$  as follows:

1. each intervened vertex  $i \in \mathcal{I}$  is split into two vertices,  $V_i(v_{\mathcal{I}})$  and  $v_i$ , and each non-intervened vertex  $j \notin \mathcal{I}$  is relabeled as  $V_j(v_{\mathcal{I}})$ .
2. the "random" vertex  $V_i(v_{\mathcal{I}})$  inherits all "incoming" edges of  $V_i$  (edges like  $* \longrightarrow V_i$  or  $* \longleftrightarrow V_i$ ) in  $G$ ;
3. the "fixed" vertex  $v_i$  inherits all "outgoing" edges of  $V_i$  (edges like  $V_i \longrightarrow *$ ) in  $G$ .

(Example on blackboard.)

## 4.4 Causal effects in linear SEM

Due to linearity, we can define the **(joint) causal effect** of  $V_{\mathcal{I}}$  on  $V$  as the matrix

$$D_{\mathcal{I}} = \{\nabla_{v_{\mathcal{I}}} V(v_{\mathcal{I}})\}^T \text{ with entries } D_{ij} = \frac{\partial}{\partial v_i} V_j(v_{\mathcal{I}}), \quad i \in \mathcal{I}, j \in [d].$$

### 4.4.1 Theorem

Suppose  $P$  is a linear SEM with respect to  $G \in \mathbb{G}(V)$  and  $\beta$  is principally stable. Then

$$\{\nabla_{v_{\mathcal{I}}} V(v_{\mathcal{I}})\}^T = \sigma(W[V_{\mathcal{I}} \rightsquigarrow V \mid V_{\mathcal{I}} \text{ in } G]).$$

Thus if  $G$  is acyclic, the total causal effect of  $V_i$  on  $V_j$  is

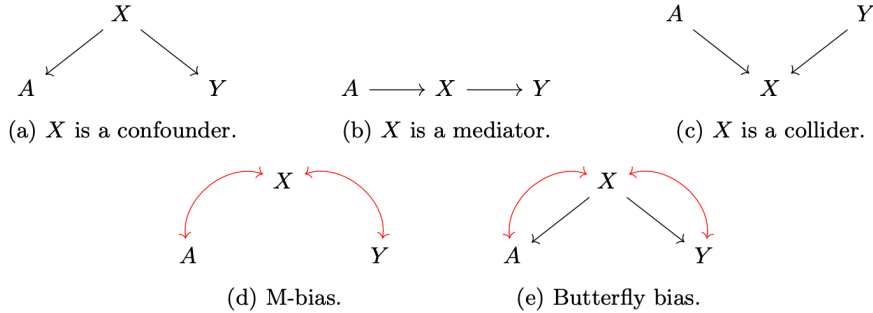
$$\frac{dV_j(v_i)}{dv_i} = \sigma(P[V_i \rightsquigarrow V_j]).$$

- Proof on blackboard.

## 4.5 Correlation is not causation

- Statistical dependence:  $\longleftrightarrow$  (local),  $\overset{t}{\longleftrightarrow}$  (global),  $t/m$ -connection (conditional).
- Causal dependence:  $\longrightarrow$  (local),  $\rightsquigarrow$  (global).

### 4.5.1 Examples



Assuming all variables are Gaussian and have unit variance, find  $\text{Cov}(A, Y)$ ,  $\text{Cov}(A, Y \mid X)$  and the causal effect of  $A$  on  $Y$ .

## 4.6 Identifiability

A central question of causal inference is identifiability. In linear models, this is asking whether the following map is injective:

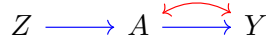
$$\Sigma : (\beta, \Lambda) \mapsto (\text{Id} - \beta)^{-T} \Lambda (\text{Id} - \beta)^{-1}.$$

We say  $\mathbb{N}^+(G)$  is

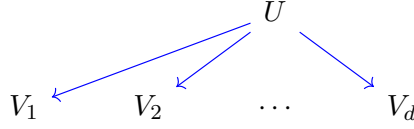
- **globally identifiable** if  $\Sigma^{-1}(\Sigma(\beta, \Lambda))$  is a singleton for all  $(\beta, \Lambda)$ ;
- **generically identifiable** if  $\Sigma^{-1}(\Sigma(\beta, \Lambda))$  is a singleton for almost all  $(\beta, \Lambda)$ ;
- **locally identifiable** if  $\Sigma^{-1}(\Sigma(\beta, \Lambda))$  does not contain an open neighborhood of  $(\beta, \Lambda)$  for all  $(\beta, \Lambda)$ .

## 4.7 Examples

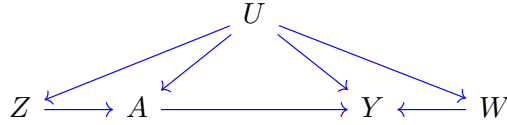
1. Instrumental variables.



2. Factor analysis ( $U$  is not observed).



3. Double negative controls ( $U$  is not observed).



## 5 Lecture 5: Conditional independence and undirected graphical models

### 5.1 Conditional independence

Let  $\mathbb{P}(\mathbb{V})$  denote all absolutely continuous probability distributions wrt  $\mathbb{V}$  (usually  $\mathbb{R}^d$ ). Let  $\mathbf{p}$  denote the density function of  $\mathbf{P} \in \mathbb{P}(\mathbb{V})$ . Consider disjoint  $V_{\mathcal{J}}, V_{\mathcal{K}}, V_{\mathcal{L}} \subseteq V$ .

#### 5.1.1 Definition

- The **conditional density function** of  $V_{\mathcal{J}}$  given  $V_{\mathcal{K}}$  is given by

$$\mathbf{p}(V_{\mathcal{J}} = v_{\mathcal{J}} \mid V_{\mathcal{K}} = v_{\mathcal{K}}) = \frac{\mathbf{p}(V_{\mathcal{J}} = v_{\mathcal{J}}, V_{\mathcal{K}} = v_{\mathcal{K}})}{\mathbf{p}(V_{\mathcal{K}} = v_{\mathcal{K}})},$$

which is well defined at any value  $v_{\mathcal{K}}$  such that  $\mathbf{p}(V_{\mathcal{K}} = v_{\mathcal{K}}) > 0$ . We often abbreviate this as  $\mathbf{p}(v_{\mathcal{J}} \mid v_{\mathcal{K}})$ .

- We write  $V_{\mathcal{J}} \perp\!\!\!\perp V_{\mathcal{K}} \mid V_{\mathcal{L}}$  **under**  $\mathbf{P}$  if one of the next equivalent conditions hold:

1.  $\mathbf{p}(v_{\mathcal{J}}, v_{\mathcal{K}} \mid v_{\mathcal{L}}) = \mathbf{p}(v_{\mathcal{J}} \mid v_{\mathcal{L}}) \mathbf{p}(v_{\mathcal{K}} \mid v_{\mathcal{L}})$ .
2.  $\mathbf{p}(v_{\mathcal{J}} \mid v_{\mathcal{L}}, v_{\mathcal{K}}) = \mathbf{p}(v_{\mathcal{J}} \mid v_{\mathcal{L}})$ .
3.  $\log \mathbf{p}(v_{\mathcal{J}}, v_{\mathcal{K}}, v_{\mathcal{L}}) = g_{\mathcal{J}, \mathcal{K}}(v_{\mathcal{J}}, v_{\mathcal{K}}) + g_{\mathcal{K}, \mathcal{L}}(v_{\mathcal{K}}, v_{\mathcal{L}})$  for some  $g_{\mathcal{J}, \mathcal{K}}$  and  $g_{\mathcal{K}, \mathcal{L}}$ .

### 5.2 Graphoid axioms

#### 5.2.1 Proposition

Consider  $\mathbf{P} \in \mathbb{P}(\mathbb{V})$  and disjoint subvectors  $J, K, L, M \subseteq V$ . We have

**Symmetry**  $(J \perp\!\!\!\perp K \mid L) \iff (K \perp\!\!\!\perp J \mid L)$ ;

**Chain rule**  $(J \perp\!\!\!\perp K \mid L, M)$  and  $(J \perp\!\!\!\perp M \mid L) \iff (J \perp\!\!\!\perp K, M \mid L)$ .

If  $\mathbf{p}(v) > 0$  for all  $v$ , then we also have

**Intersection**  $(J \perp\!\!\!\perp K \mid L, M)$  and  $(J \perp\!\!\!\perp M \mid K, L) \implies (J \perp\!\!\!\perp K, M \mid L)$ .

(Proof is left as an exercise.)

- A ternary relation that satisfy these axioms is called a **graphoid**. The terminology is justified by the following visualization.

|     |     |     |
|-----|-----|-----|
| $J$ | $L$ | $K$ |
|     |     | $M$ |

### 5.3 Separation in undirected graphs

#### 5.3.1 Definition

- Let  $\mathbb{UG}(V)$  denote the collection of all simple undirected graphs with vertex set  $V$ .
- Given  $G \in \mathbb{UG}(V)$  and disjoint subsets  $J, K, L \subset V$ , we say  $L$  **separate**  $J$  and  $K$  in  $G$  and write

$$\text{not } J \text{ --- } * \text{ --- } K \mid L \text{ in } G$$

if every path from a vertex in  $J$  to a vertex in  $K$  in  $G$  contains a non-endpoint in  $L$ .

#### 5.3.2 Interpretation

This is "dual" to separation in bidirected graphs:

- $J \text{ --- } L \text{ --- } K, J \longleftrightarrow M \longleftrightarrow K$  are blocked given  $L$ ;
- $J \text{ --- } M \text{ --- } K, J \longleftrightarrow L \longleftrightarrow K$  are not blocked given  $L$ .

### 5.4 Undirected graphical models

- Let  $\mathbb{P}_{\text{GM}}(G)$  collects all  $P \in \mathbb{P}(\mathbb{V})$  that satisfies the **global Markov property** wrt  $G \in \mathbb{UG}(V)$ : for all disjoint  $J, K, L \subset V$ ,

$$\text{not } J \text{ --- } * \text{ --- } K \mid L \text{ in } G \implies J \perp\!\!\!\perp K \mid L \text{ under } P.$$

- Let  $\mathbb{P}_{\text{F}}(G)$  collects all  $P \in \mathbb{P}(\mathbb{V})$  that **factorizes** wrt  $G \in \mathbb{UG}(V)$ :

$$p(v) = \prod_{V_{\mathcal{J}} \in \mathcal{C}(G)} f_{\mathcal{J}}(v_{\mathcal{J}}),$$

for some  $f_{\mathcal{J}}, \mathcal{J} \subseteq [d]$ , where  $\mathcal{C}(G)$  collects all "cliques" (complete subgraphs) of  $G$ .

#### 5.4.1 Theorem (Hammersley-Clifford)

For any  $G \in \mathbb{UG}(V)$ , we have  $\mathbb{P}_{\text{F}}(G) \subseteq \mathbb{P}_{\text{GM}}(G)$  and  $\mathbb{P}_{\text{F}}^+(G) = \mathbb{P}_{\text{GM}}^+(G)$ .

- Proof of the first part. ( $\mathbb{P}^+$  means positive density functions.)

### 5.5 Graph augmentation

Graph separations in undirected and directed graphs are closely related via the augmentation map  $\text{aug} : \mathbb{G}^*(V) \rightarrow \mathbb{UG}(V)$  defined by

$$V_j \text{ --- } V_k \text{ in } \text{aug}(G) \iff V_j \longleftrightarrow * \longleftrightarrow V_k \text{ in } G, \text{ for all } V_j \neq V_k.$$

- When restricted to  $\mathbb{G}_{\text{DA}}^*(V)$ , this is called **moralization** in the literature because it connects any two parents of the same child.
- For  $J \subseteq V$ , define  $\text{an}(J) = \{V_k \in V : V_k \rightsquigarrow J \text{ in } G\}$  and  $\overline{\text{an}}(J) = \text{an}(J) \cup J$  (the smallest ancestral set containing  $J$ ).

### 5.5.1 Proposition

For any  $G \in \mathbb{G}^*(V)$  and disjoint  $J, K, L \subset V$ , we have, with  $\tilde{V} = \overline{\text{an}}(J \cup K \cup L)$ ,

$$J \rightsquigarrow * \rightsquigarrow K \mid L \text{ in } G \iff J \text{ --- } * \text{ --- } K \mid L \text{ in } \text{aug} \circ \text{margin}_{\tilde{V}}(G).$$

- Proof on blackboard.

## 6 Lecture 6: DAG models and ADMG models

### 6.1 DAG models

#### 6.1.1 Definition

- Let  $\mathbb{P}_F(G)$  collects all  $P \in \mathbb{P}(V)$  that **factorizes** wrt  $G \in \mathbb{G}_{DA}^*(V)$ :

$$p(v) = \prod_{j=1}^p p(v_j \mid v_{\text{pa}_G(j)}).$$

- Let  $\mathbb{P}_{GM}(G)$  collects all  $P \in \mathbb{P}(V)$  that satisfies the **global Markov property** wrt  $G \in \mathbb{G}_{DA}^*(V)$ : for any disjoint subsets  $J, K, L \subset V$ ,

$$\text{not } J \overset{d}{\rightsquigarrow} * \overset{d}{\rightsquigarrow} K \mid L \text{ in } G \implies J \perp\!\!\!\perp K \mid L \text{ under } P.$$

#### 6.1.2 Theorem

For any  $G \in \mathbb{G}_{DA}^*(V)$ , we have  $\mathbb{P}_F(G) = \mathbb{P}_{GM}(G)$ .

- Proof on blackboard.

### 6.2 ADMG models

- Let  $\mathbb{P}_{GM}(G)$  collects all  $P \in \mathbb{P}(V)$  that satisfies the **global Markov property** wrt  $G \in \mathbb{G}_A^*(V)$ : for any disjoint subsets  $J, K, L \subset V$ ,

$$\text{not } J \rightsquigarrow * \rightsquigarrow K \mid L \text{ in } G \implies J \perp\!\!\!\perp K \mid L \text{ under } P.$$

- Alternatively, we can define ADMG models by using simpler expanded graphs.

#### 6.2.1 Graph expansion

- We say  $G'$  is an **expansion** of  $G \in \mathbb{G}^*(V)$  if it is in

$$\text{expand}(G) = \text{margin}_V^{-1}(G) = \bigcup_{V' \supseteq V} \{G' \in \mathbb{G}^*(V') : \text{margin}_V(G') = G\}.$$

- Often, bidirected edges in  $G$  correspond to certain latent variables in  $G'$ .
- If we are satisfied with an expansion  $G'$  of  $G$ , we can use  $\text{margin}_V(\mathbb{P}_{GM}(G'))$  as the model for  $G$ .

### 6.3 ADMG models

- **Pairwise expansion:** replace every  $V_j \longleftrightarrow V_k$  with  $V_j \xleftarrow{U_{jk}} U_{jk} \xrightarrow{V_k} V_k$ .
- **Clique expansion:** replace every bidirected clique  $V_{\mathcal{J}}$  (complete bidirected subgraph) with directed edges  $U_{\mathcal{J}} \rightarrow V_j, j \in \mathcal{J}$ .
- **Noise expansion:** add  $U_j \rightarrow V_j$  such that  $U_j$  inherits all bidirected edges of  $V_j$ .
- Example on blackboard.

Let the corresponding models be denoted as  $\mathbb{P}_{PE}(G)$ ,  $\mathbb{P}_{CE}(G)$ ,  $\mathbb{P}_{NE}(G)$ .

### 6.3.1 Proposition

For any  $G \in \mathbb{G}_A^*(V)$ , we have

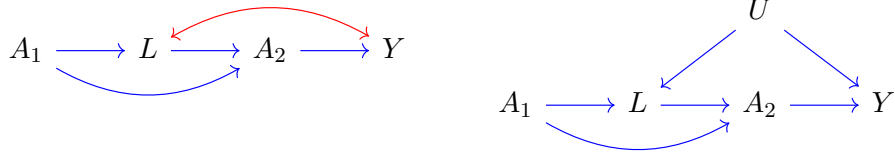
$$\mathbb{P}_{PE}(G) \subseteq \mathbb{P}_{CE}(G) \subseteq \mathbb{P}_{NE}(G) \subseteq \mathbb{P}_{GM}(G),$$

and  $\subseteq$  in above cannot be replaced by  $=$  in general.

- The latent variable models (PE/CE/NE) have additional equality and inequality constraints.

### 6.4 Additional equality constraints

Consider the following (trimmed) ADMG and its expansion ( $U$  is latent).



Suppose  $P \in \mathbb{P}_{PE}(G)$  (PE,CE,NE are actually equivalent in this example). Then

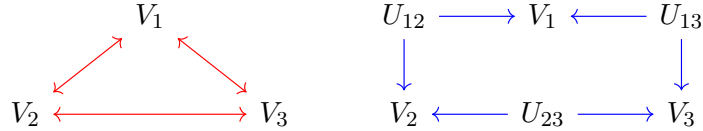
$$\int p(y \mid a_1, l, a_2) p(l \mid a_1) dl \text{ does not depend on } a_1.$$

- Proof on blackboard.
- This can be understood as a "hidden" independence  $Y \perp\!\!\!\perp A_1$  in the kernel

$$p(a_1, l, y \mid \text{do}(a_2)) = \frac{p(a_1, l, a_2, y)}{p(a_2 \mid a_1, l)} = p(a_1) p(l \mid a_1) p(y \mid a_1, l, a_2).$$

### 6.5 Additional inequality constraints

Consider the following bidirected clique and its pairwise expansion.



If  $P \in \mathbb{P}_{PE}(G)$  and  $\mathbb{V} = \{-1, 1\}^3$ , the following "perfect correlation" is impossible:

$$P(V_1 = V_2 = V_3 = 1) = P(V_1 = V_2 = V_3 = -1) = \frac{1}{2}.$$

- Heuristically, if  $V_1 = V_2$ , then  $V_1$  cannot depend on  $U_{13}$ .
- $\mathbb{P}_{CE}(G)$  or  $\mathbb{P}_{NE}(G)$  have no such constraints.
- Other related examples: Bell's inequality in quantum mechanics; Balke-Pearl bound for instrumental variable graph.

## 6.6 Advantages of the noise expansion model

1. It is equivalent to a **natural nonparametric generalization** of the linear SEM: if  $P \in \mathbb{P}_{\text{NE}}(G)$ , then  $V$  satisfies

$$V_j = f_j(V_{\text{pa}_G(j)}, E_j)$$

for some functions  $f_1, \dots, f_d$  and noise variables  $E_1, \dots, E_d$  that satisfy

$$V_{\mathcal{J}} \not\leftrightarrow V_{\mathcal{K}} \text{ in } G \implies E_{\mathcal{J}} \perp\!\!\!\perp E_{\mathcal{K}}.$$

2. Let  $\mathbb{G}_{\text{UA}}^*(V)$  collect all **unconfounded ADMGs** ( $V_j \leftrightarrow V_k$  in  $G$ ,  $V_j \neq V_k$  implies  $\text{pa}_G(j) = \emptyset$ ) with vertex set  $V$ . Then

- For all  $G \in \mathbb{G}_{\text{UA}}^*(V)$ , we have  $\mathbb{P}_{\text{NE}}(G) = \mathbb{P}_{\text{GM}}(G)$ .
- For all  $G \in \mathbb{G}_{\text{A}}^*(V)$ , we have

$$\mathbb{P}_{\text{NE}}(G) = \bigcup_{V' \supseteq V} \bigcup_{G'} \text{margin}_V(\mathbb{P}_{\text{NE}}(G')).$$

(The second union is over  $G' \in \text{expand}(G) \cap \mathbb{G}_{\text{UA}}^*(V')$ .)

## 7 Lecture 7: Causal Markov model

### 7.1 Causal Markov model

Recall that a causal model is a collection of consistent probability distributions on the potential outcomes schedule  $V(\cdot)$ .

#### 7.1.1 Definition

Let  $\mathbb{CP}(G)$  collect all distributions  $P$  on  $V(\cdot)$  that is **causal Markov** wrt  $G \in \mathbb{G}_{\text{A}}^*(V)$ :

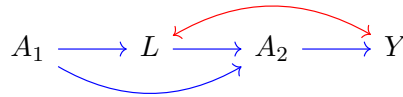
1. Recursive substitution: With  $P$ -probability 1, we have

$$V_j(v_{\mathcal{I}}) = V_j(v_{\text{pa}_G(j) \cap \mathcal{I}}, V_{\text{pa}_G(j) \setminus \mathcal{I}}(v_{\mathcal{I}})) \text{ for all } j \in [d], \mathcal{I} \subseteq [d], v \in \mathbb{V}.$$

2. Basic potential outcomes are Markov wrt bidirected subgraph:

$$V_{\mathcal{J}} \not\leftrightarrow V_{\mathcal{K}} \text{ in } G \implies V_{\mathcal{J}}(v) \perp\!\!\!\perp V_{\mathcal{K}}(v) \text{ under } P \text{ for all disjoint } \mathcal{J}, \mathcal{K} \subset [d] \text{ and } v \in \mathbb{V}.$$

Example: What is  $Y(a_1, a_2)$ ? What is  $Y(a_1)$ ?



### 7.2 Properties

Suppose  $G \in \mathbb{G}_{\text{A}}^*(V)$  and  $P \in \mathbb{CP}(G)$ .

#### 7.2.1 Property 1 (Consistency of potential outcomes)

$P(V(v_{\mathcal{I}}, v_{\mathcal{J}}) = V(v_{\mathcal{I}}) \mid V_{\mathcal{J}}(v_{\mathcal{I}}) = v_{\mathcal{J}}) = 1$ , for all disjoint  $\mathcal{I}, \mathcal{J} \subseteq [d]$ ,  $v \in \mathbb{V}$ .

#### 7.2.2 Property 2 (Simplifying potential outcomes)

For any  $V_{\mathcal{J}}, V_{\mathcal{K}}, V_{\mathcal{L}} \subseteq V$ ,  $V_{\mathcal{K}} \cap V_{\mathcal{L}} = \emptyset$ , we have

$$\text{not } V_{\mathcal{L}} \rightsquigarrow V_{\mathcal{J}} \mid V_{\mathcal{K}} \text{ in } G \implies P(V_{\mathcal{J}}(v_{\mathcal{K}}, v_{\mathcal{L}}) = V_{\mathcal{J}}(v_{\mathcal{K}})) = 1, \text{ for all } v_{\mathcal{K}} \in \mathbb{V}_{\mathcal{K}}, v_{\mathcal{L}} \in \mathbb{V}_{\mathcal{L}}.$$

### 7.2.3 Property 3 (SWIG Markov property)

We have  $\text{margin}_{V(v_I)}(\mathbf{P}) \in \mathbb{P}_{\text{GM}}(\mathbf{G}(v_I))$  for all  $V_I \subseteq V$  and  $v \in \mathbb{V}$ .

- Proof on blackboard.

## 8 Lecture 8: Causal identification and confounder selection

### 8.1 Identification by fixing

Consider  $\mathbf{G} \in \mathbb{G}_A^*(V)$ .

- We say  $V_j \in V$  is **fixable** in  $\mathbf{G}$  if there exists no  $V_k \in V$  such that  $V_j \rightsquigarrow V_k$  and  $V_j \longleftrightarrow * \longleftrightarrow V_k$  in  $\mathbf{G}$ .
- For  $V_j \in V$ , its **Markov background** in  $\mathbf{G}$  is defined as

$$\text{mbg}_{\mathbf{G}}(V_j) = \{V_k \in V : V_k \longleftrightarrow * \longleftrightarrow V_j \text{ in } \mathbf{G}\}.$$

#### 8.1.1 Proposition

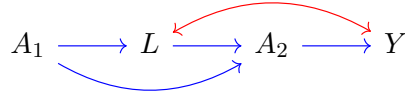
Consider  $\mathbf{G} \in \mathbb{G}_A^*(V)$  and  $\mathbf{P} \in \mathbb{CP}(\mathbf{G})$ . If  $V_j \in V$  is fixable in  $\mathbf{G}$ , then

$$\frac{\mathbf{p}(V_j(v_j) = \tilde{v}_j, V_{-j}(v_j) = v_{-j})}{\mathbf{p}(V_j = v_j, V_{-j} = v_{-j})} = \frac{\mathbf{p}(V_j = \tilde{v}_j \mid V_{\text{mbg}(j)} = v_{\text{mbg}(j)})}{\mathbf{p}(V_j = v_j \mid V_{\text{mbg}(j)} = v_{\text{mbg}(j)})}, \text{ for all } v \in \mathbb{V} \text{ and } v_j^* \in \mathbb{V}_j,$$

whenever  $\mathbf{p}(V_j = v_j \mid V_{\text{mbg}(j)} = v_{\text{mbg}(j)}) > 0$ .

- Proof on blackboard.

### 8.2 Example



Show that the equality constraint

$$\int \mathbf{p}(y \mid a_1, l, a_2) \mathbf{p}(l \mid a_1) dl \text{ does not depend on } a_1.$$

corresponds to

- the independence  $Y(a_2) \perp\!\!\!\perp A_1$ ; or
- no direct  $A_1 \rightarrow Y$  effect:  $Y(a_1, a_2) = Y(a_2)$ .

### 8.3 Back-door criterion

Consider  $\mathbf{G} \in \mathbb{G}_A^*(V)$ ,  $\mathbf{P} \in \mathbb{CP}(\mathbf{G})$ ,  $A, Y \in V$ . Interested in the causal effect of  $A$  on  $Y$ .

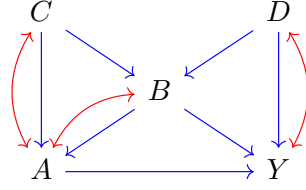
#### 8.3.1 Theorem

Suppose  $X \subset V$ ,  $X \cap \{A, Y\} = \emptyset$  satisfies

1.  $A \not\rightsquigarrow X$  in  $\mathbf{G}$ ;
2.  $P[A \longleftrightarrow * \longleftrightarrow Y \mid_a X] = \emptyset$ .

Then  $\mathbf{p}(Y(a) = y \mid X = x) = \mathbf{p}(Y = y \mid A = a, X = x)$ .

- Proof on blackboard.
- Example: Which  $X \subseteq \{B, C, D\}$  meet the back-door criterion?



## 8.4 Confounder selection

Can we select a set of confounders  $X$  without knowing the full graph  $G$ ?

### 8.4.1 Definition (symmetric back-door criterion)

- $X \subseteq V \setminus \{A, Y\}$  is an **adjustment set** for  $A, Y \in V$  if  $A \not\rightsquigarrow X$  and  $Y \not\rightsquigarrow X$ .
- An adjustment set  $X$  is **sufficient** if  $P[A \longleftrightarrow * \longleftrightarrow Y \mid_a X] = \emptyset$ .
- An adjustment set  $X$  is **primary** if  $P[A \longleftrightarrow Y \mid X] = \emptyset$ .

### 8.4.2 Heuristics

Directly blocking all confounding paths is difficult, because

$$P[A \longleftrightarrow * \longleftrightarrow Y \mid_a X] = \emptyset \not\Rightarrow P[A \longleftrightarrow * \longleftrightarrow Y \mid_a \tilde{X}] = \emptyset \text{ for } X \subset \tilde{X}.$$

But we can block confounding arcs recursively, because

$$P[A \longleftrightarrow Y \mid X] = \emptyset \Rightarrow P[A \longleftrightarrow Y \mid X'] = \emptyset \text{ for } X \subset X'.$$

## 8.5 District criterion

### 8.5.1 Theorem (marginalization preserves confounding arcs and paths)

Consider  $G \in \mathbb{G}_A^*(V)$ , distinct  $A, Y \in V$ ,  $X \subseteq V \setminus \{A, Y\}$ . For any vertex set  $\tilde{V}$  such that  $\{A, B\} \cup C \subseteq \tilde{V} \subseteq V$ , we have

$$\begin{aligned} P[A \longleftrightarrow Y \mid X \text{ in } G] = \emptyset &\iff P[A \longleftrightarrow Y \mid X \text{ in } \text{margin}_{\tilde{V}}(G)] = \emptyset, \\ P[A \longleftrightarrow * \longleftrightarrow Y \mid_a X \text{ in } G] = \emptyset &\iff P[A \longleftrightarrow * \longleftrightarrow Y \mid_a X \text{ in } \text{margin}_{\tilde{V}}(G)] = \emptyset. \end{aligned}$$

As a corollary, we have

$$\begin{aligned} P[A \longleftrightarrow Y \mid X \text{ in } G] = \emptyset &\iff \text{not } A \longleftrightarrow Y \text{ in } \text{margin}_{\{A, Y\} \cup X}(G), \\ P[A \longleftrightarrow * \longleftrightarrow Y \mid_a X \text{ in } G] = \emptyset &\iff \text{not } A \longleftrightarrow * \longleftrightarrow Y \text{ in } \text{margin}_{\{A, Y\} \cup X}(G). \end{aligned}$$

## 8.6 Iterative graph expansion

```

1: procedure CONFOUNDERSELECT( $A, Y$ )
2:    $\mathcal{R} = \emptyset$ 
3:   procedure GRAPHEXPAND( $X, \mathcal{B}_y, \mathcal{B}_n$ )
4:     if  $A \longleftrightarrow * \longleftrightarrow Y$  by edges in  $\mathcal{B}_y$  then
5:       return
6:     else if not  $A \longleftrightarrow * \longleftrightarrow Y$  by edges in  $(X \cup \{A, Y\}) \times (X \cup \{A, Y\}) \setminus \mathcal{B}_n$  then
7:        $\mathcal{R} = \mathcal{R} \cup \{X\}$ 
8:       return
9:     end if
10:     $(C \leftrightarrow D) = \text{SELECTEDGE}(A, Y, X, \mathcal{B}_y, \mathcal{B}_n)$ 
11:    for  $X'$  in  $\text{FINDPRIMARY}(C \leftrightarrow D, X)$  do

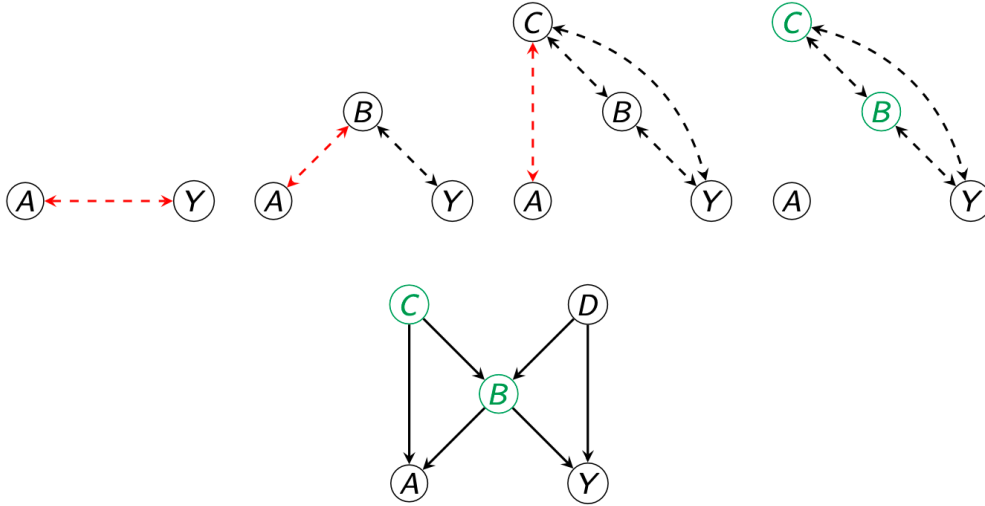
```

```

12:     GRAPHEXPAND( $X \cup X'$ ,  $\mathcal{B}_y$ ,  $\mathcal{B}_n \cup \{C \leftrightarrow D\}$ )
13:   end for
14:   GRAPHEXPAND( $X$ ,  $\mathcal{B}_y \cup \{C \leftrightarrow D\}$ ,  $\mathcal{B}_n$ )
15: end procedure
16: GRAPHEXPAND( $\emptyset$ ,  $\emptyset$ ,  $\emptyset$ )
17: return  $\mathcal{R}$ 
18: end procedure

```

## 8.7 Illustration



Shiny app: <https://ricguo.shinyapps.io/InteractiveConfSel/>